# MiniPascal Compiler - Final Project Report (Code Generation Phase)

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### 1. Introduction & Project Conclusion

This report details the third and final phase of the MiniPascal Compiler project: **Code Generation**. Following the successful completion of the frontend (Lexical and Syntax Analysis) and the middle-end (Semantic Analysis), this phase completes the compiler's primary objective: to translate a semantically-valid Abstract Syntax Tree (AST) into executable assembly code for our custom, stack-based Virtual Machine.

With the implementation of the code generator, the project is now **functionally complete**. The compiler can successfully take a MiniPascal source file, process it through all compilation stages, and produce a .vm file that can be executed by the target VM. This report outlines the architecture and capabilities of the code generation backend.

### 2. Core Technologies & Tools Update

The core technologies remain consistent, with the addition of the final backend component:

* **Lexical Analysis (Scanner):** GNU Flex
* **Syntax Analysis (Parser):** GNU Bison
* **Abstract Syntax Tree (AST) & Semantic Analysis:** C++ (std=c++17)
* **Code Generation (Backend):** A new CodeGenerator class (codegenerator.h, codegenerator.cpp), implementing the SemanticVisitor pattern, is responsible for traversing the AST and emitting assembly code.
* **Build System:** GNU Make with g++ compiler
* **Development Environment:** MinGW64 on Windows

### 3. Code Generation (The Backend)

The code generator is the final stage of the compiler. It traverses the AST, which has been validated and annotated by the semantic analyzer, and produces a sequence of text-based assembly instructions compatible with our target Virtual Machine.

#### Key Capabilities & Translation Strategies:

**3.1. Stack Frame & Scope Management**

The generator correctly manages the runtime environment for global and local scopes.

* **Global Scope:** Global variables are allocated at the start of the program using pushn and accessed directly via their offsets from the gp (global pointer) register using pushg and storeg.
* **Subprogram Scope:**
  + At the beginning of a subprogram, a single pushn instruction allocates all necessary space for its local variables on the stack.
  + Local variables are accessed via positive offsets from the fp (frame pointer) register using pushl and storel.
  + Parameters passed to the subprogram are accessed via negative offsets from the fp register (e.g., pushl -1, pushl -2).
  + The call and return VM instructions are used to manage the call stack, saving and restoring the pc (program counter) and fp.

**3.2. Expression Translation**

Expressions are translated by recursively visiting the expression tree and emitting instructions that place operands on the stack before the operator.

* **Arithmetic Operations:** +, -, \*, div are translated to add, sub, mul, div. The real division / is translated to fdiv.
* **Type Promotion:** The generator correctly emits the itof instruction when an INTEGER is used in a REAL context (e.g., 5 + 2.5).
* **Relational** & **Logical Operations:** Relational operators (=, >, <) are translated to equal, sup, inf, etc. Logical operators (AND, OR, NOT) are translated into corresponding arithmetic and comparison sequences.

**3.3. Statement & Control Flow Translation**

Pascal statements are translated into labeled blocks of assembly code.

* **IF...THEN...ELSE:** Implemented using a conditional jump (jz). The condition is evaluated, and jz is used to skip to the ELSE block's label if the result is false. A jump instruction ensures the THEN block skips over the ELSE block.
* **WHILE...DO:** Implemented with a starting label for the condition check and an ending label. A jz instruction exits the loop, and an unconditional jump at the end of the body returns to the condition check.

**3.4. Subprogram Calling Convention**

A standard calling convention has been implemented for both procedures and functions.

* **Caller Responsibilities:**
  1. For functions, the caller first allocates space for the return value (pushn 1).
  2. Pushes all arguments onto the stack in **right-to-left** order.
  3. Pushes the target subprogram's address (pusha SubprogramName).
  4. Executes the call instruction.
  5. After the subprogram returns, it cleans up the arguments from the stack (pop N).
* **Callee Responsibilities:**
  1. The subprogram's code begins execution.
  2. It allocates space for its local variables.
  3. When a RETURN statement is encountered in a function, the result is calculated and stored in the pre-allocated return value slot [fp - (num\_params + 1)] using storel.
  4. The return instruction restores the caller's stack frame and transfers control.

**3.5. Built-in I/O Procedures**

The generator has specific logic to handle the built-in I/O procedures.

* **write, writeln:** Arguments are evaluated one by one. Depending on the argument's type, the appropriate VM instruction (writei, writef, writes) is emitted. writeln additionally prints a newline character.
* **read, readln:** For each variable argument, the generator emits read (to get a string from the user) followed by atoi or atof to convert it to the variable's type. The final result is stored using storeg or storel.

### 4. Final Project Status

The MiniPascal compiler project is now **functionally complete**. It successfully implements the entire compilation pipeline for the specified language subset.

* **Full Language Support:** All features defined in the MiniPascal specification—including global/local variables, all data types, arrays, IF and WHILE statements, and full subprogram support with recursion—are parsed, semantically analyzed, and compiled into executable code.
* **Verified Correctness:** A comprehensive test suite of over 30 focused tests has been used to validate the compiler's output and ensure its correctness across a wide range of features and edge cases.
* **Executable Output:** The primary goal has been achieved: the compiler produces .vm assembly files that run correctly on the target virtual machine.

**Known Issues:**

* There is a known bug causing a "VM error: Illegal Operand" when accessing **local arrays** declared inside a procedure or function. All other features, including global arrays and local simple variables, are working correctly.

### 5. Illustrative Example: Final Output

The following MiniPascal code demonstrates a combination of features including function calls, local variables, and I/O.

**Source Code (Test00.pas can be found in Tests/Test00.pas):**

PROGRAM FullTest;  
VAR  
 x, y, z: INTEGER;  
   
FUNCTION addition(a: INTEGER; b: INTEGER): INTEGER;  
 VAR  
 result: INTEGER;  
 BEGIN  
 result := a + b;  
 RETURN result;  
 END;  
  
BEGIN  
 x := 15;  
 y := 27;  
 z := addition(x, y);  
 write('The result of add(15, 27) is: ', z);  
 writeln;  
END.

**Generated Assembly (output.vm):**

start  
 jump main\_entry  
 jump addition\_end  
addition:  
 pushn 1  
 pushl -1  
 pushl -2  
 add  
 storel 0  
 pushl 0  
 storel -3  
 return  
addition\_end:  
main\_entry:  
 pushn 3  
 pushi 15  
 storeg 0  
 pushi 27  
 storeg 1  
 pushn 1  
 pushg 1  
 pushg 0  
 pusha add  
 call  
 pop 2  
 storeg 2  
 pushs "The result of add(15, 27) is: "  
 writes  
 pushg 2  
 writei  
 pushs "\n"  
 writes  
 stop